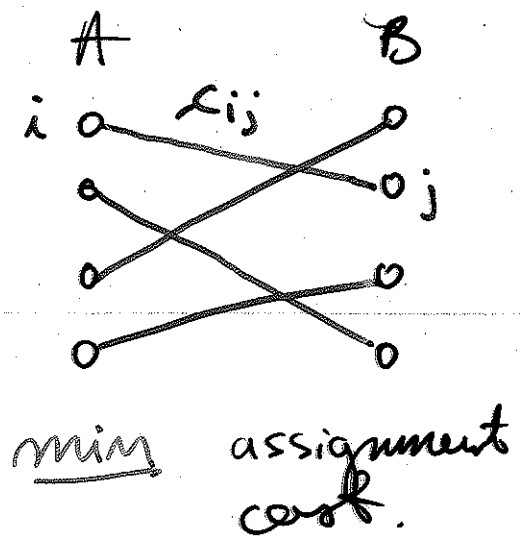


# Assignment

$$(P) \begin{cases} \min \sum_{\substack{i \in A \\ j \in B}} c_{ij} x_{ij} \\ \sum_{j \in B} x_{ij} = 1, \quad i \in A \\ \sum_{i \in A} x_{ij} = 1, \quad j \in B \\ x_{ij} \geq 0 \end{cases}$$



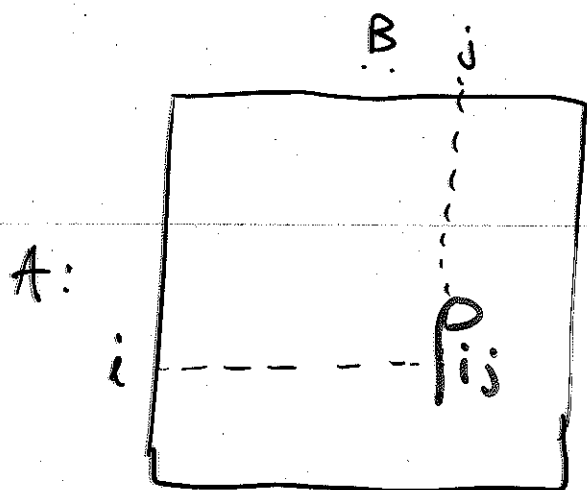
$$(D) \begin{cases} \max \sum_{i \in A} \alpha_i + \sum_{j \in B} \beta_j \\ \alpha_i + \beta_j \leq c_{ij}, \quad \substack{i \in A \\ j \in B} \end{cases}$$

Reduced cost (slack of dual constraint)

$$P_{ij} = c_{ij} - \alpha_i - \beta_j$$

$$C.S.: \quad x_{ij} \cdot P_{ij} = 0.$$

# Primal-dual algorithm



- dual feasibility

$$P_{ij} \geq 0$$

- C.S.:

assign  $i$  to  $j$   
( $x_{ij} = 1$ ) only if

$$P_{ij} = 0$$

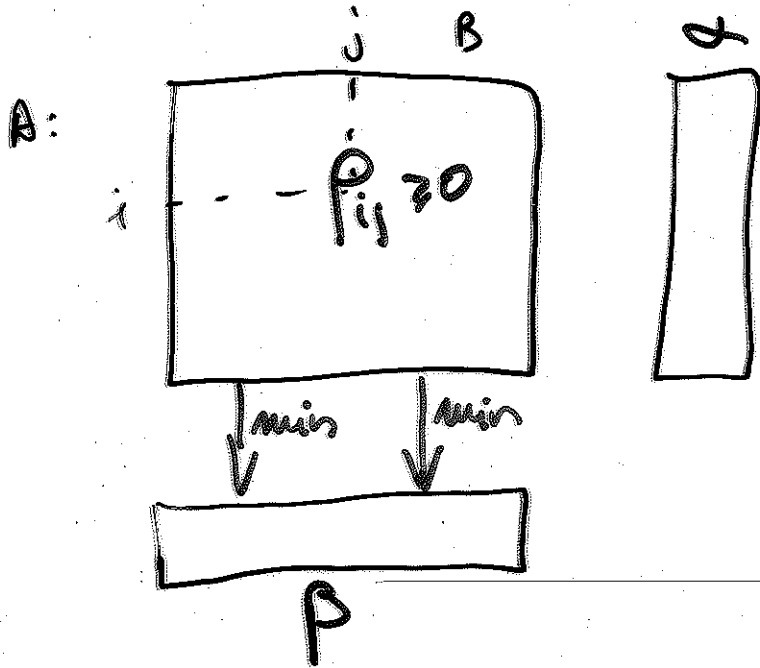
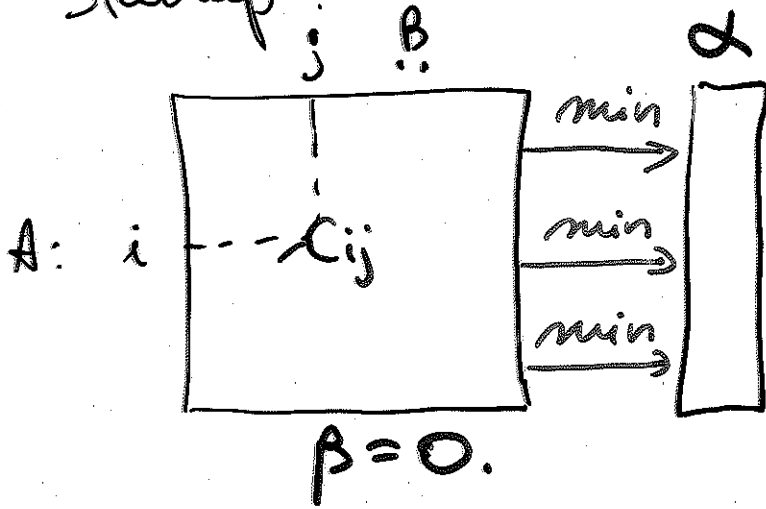
maintained throughout alg.

- primal feasibility:

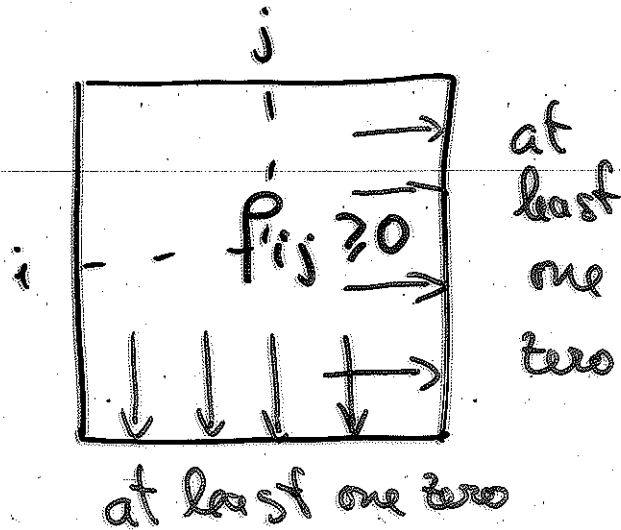
attempted through combinatorial methods (select 0 cells for assign.)

# Primal-dual (c'ed).

Startup:



Result:

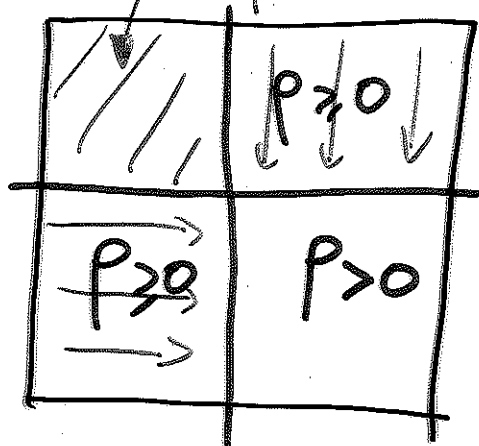


# Primal-dual (c'ed)

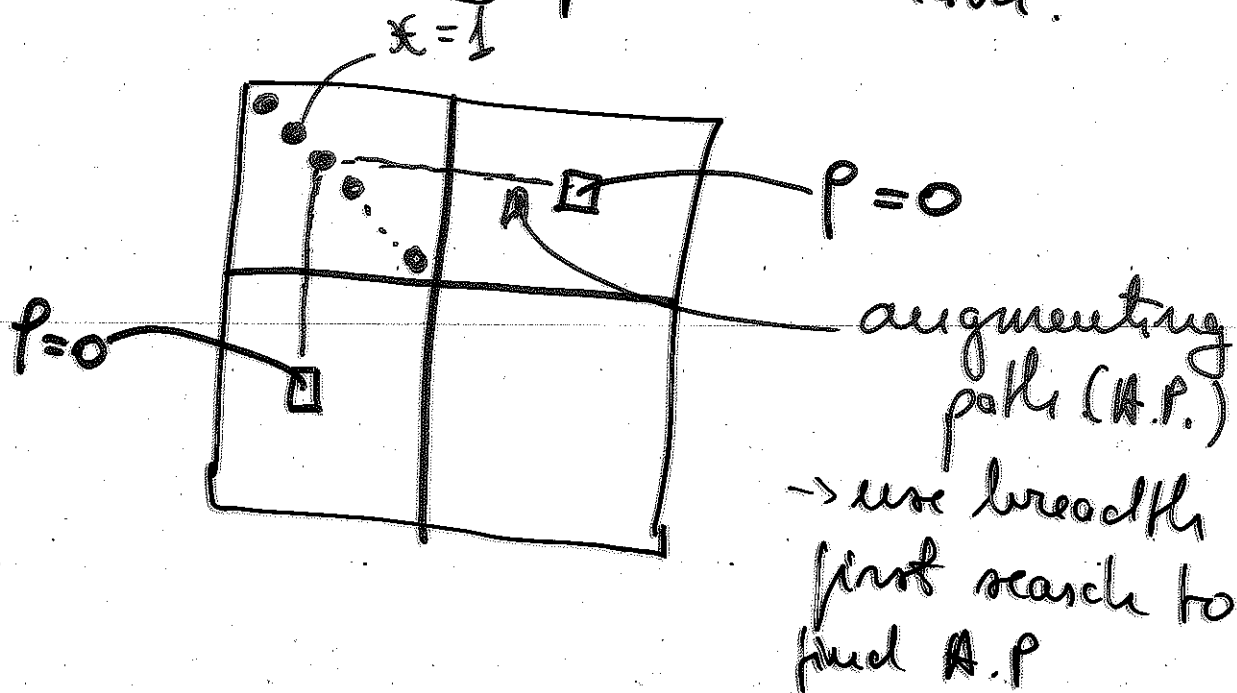
iteration: dual variables fixed ( $p_{ij}$ ),  
increase primal feasibility.

a) Greedy algorithm { - select cell  $p_{ij} = 0$   
- assign  $x_{ij} = 1$  if no conflicts

b)

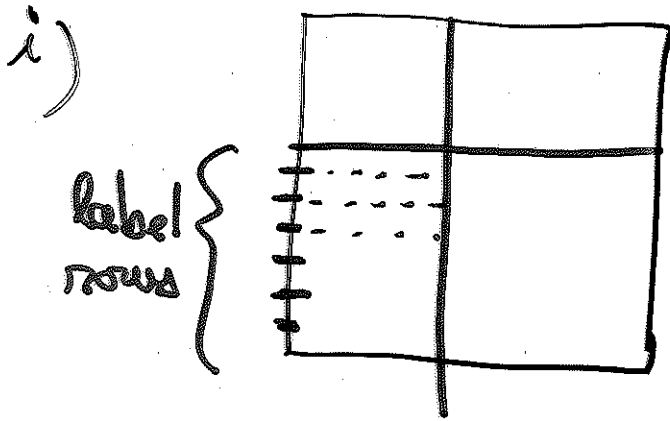


- extending primal solution:

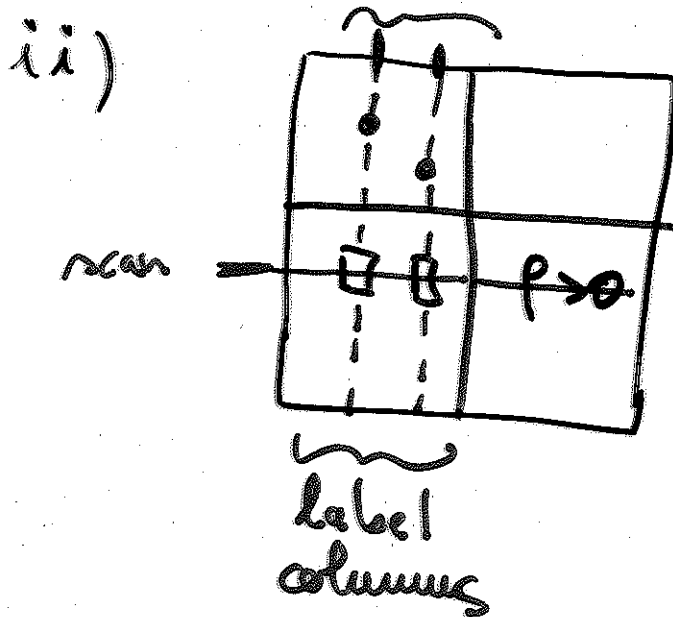


Primal-dual (c'ed).

c) Breadth first search.



- add labelled rows to queue.



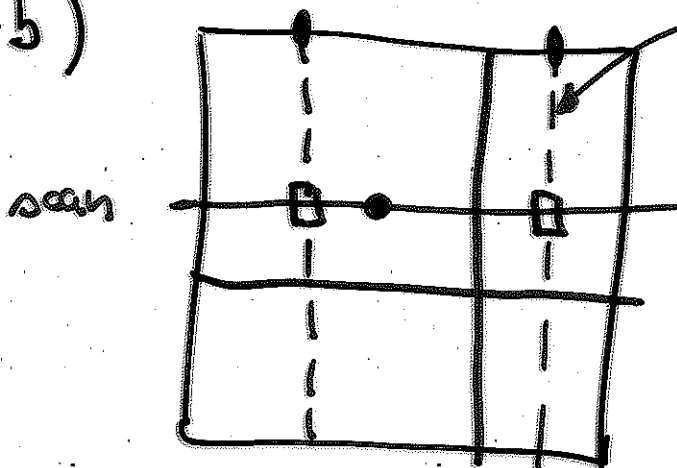
- row scan

$\square: p=0$

Label cols with  $p=0$  (if unlabelled)

Add to queue

ii-b)



no allocation

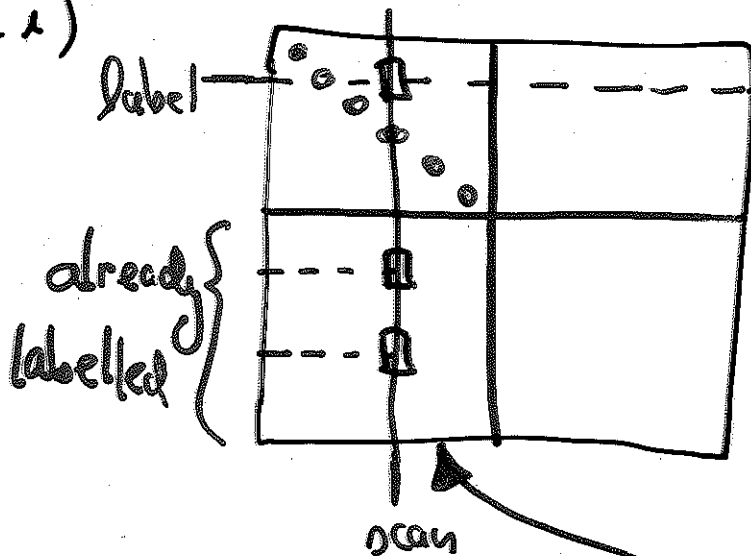
- row scan:

If a column is labeled & has no allocation,

AP found!

# c) Breadth first search (c'ed)

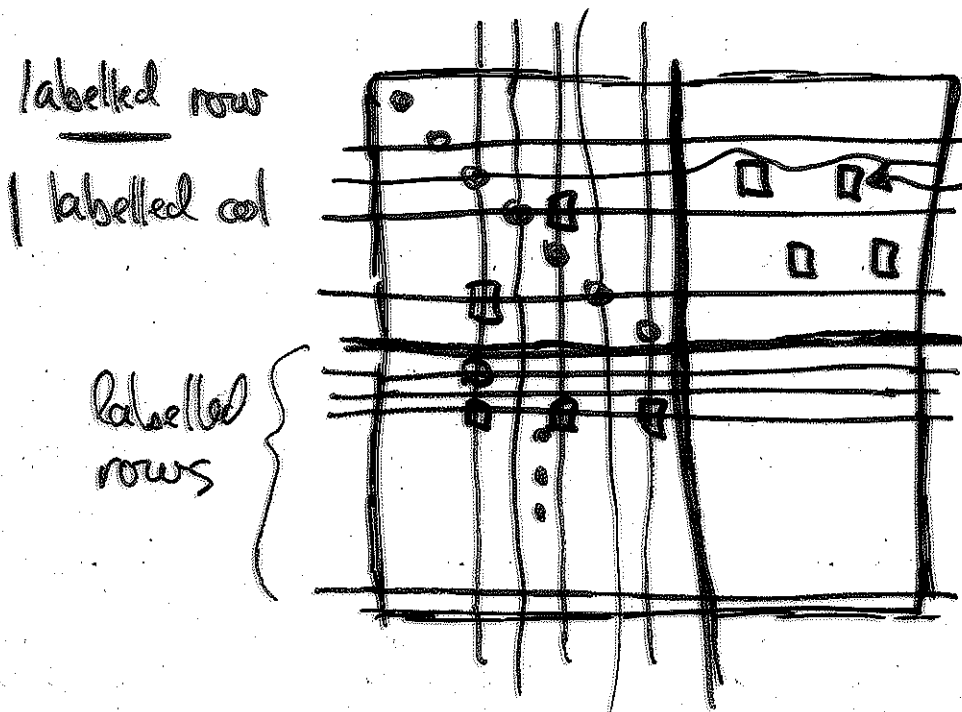
iii)



- Column scan:  
Label unlabelled rows;  
Add them to queue.

Labelled rows & cols = rows & cols reachable from region by following row-col paths through  $(P=0)$  cells.

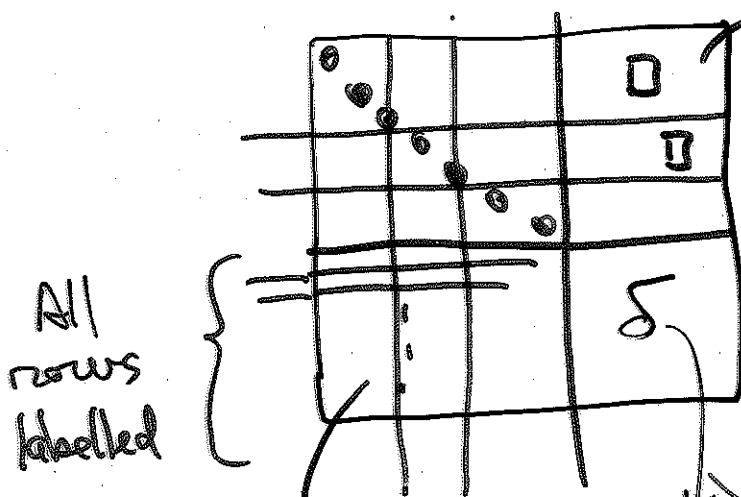
No A.P.



All  $P=0$  cells are NOT on labelled rows (otherwise A.P.!) )

# Primal-dual (c'ed)

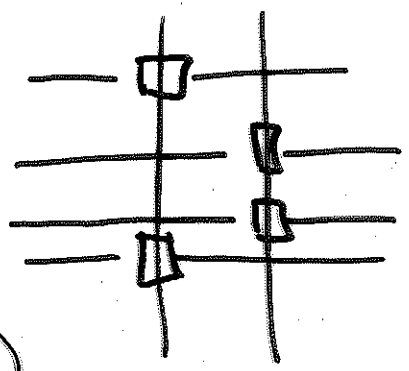
d) Change dual:

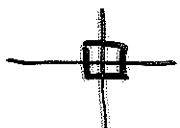


No  $P=0$  cells on labelled rows

All rows labelled

incorrect

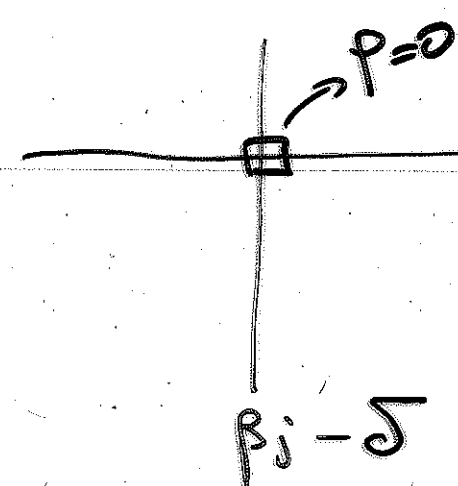


All  $P=0$  cells on labelled rows ~~not~~ sit on lab. cols. 

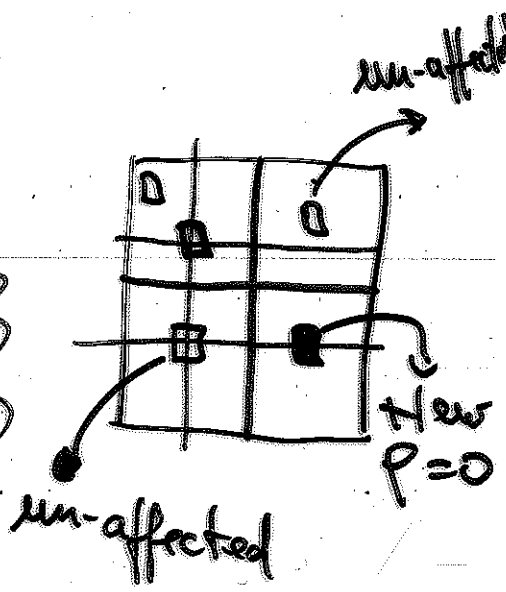
~~$\delta = \min P_{ij}$~~

All  $P=0$  cells of labelled cols sit on labelled rows bc, lab. cols = scanned!

New duals:

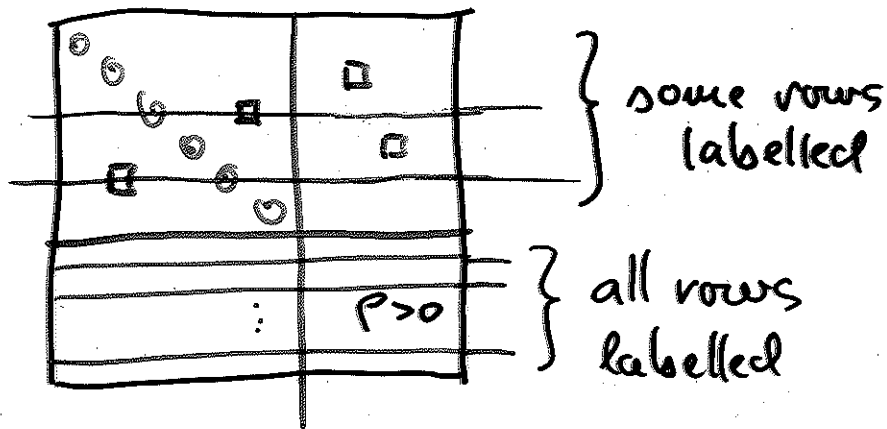


The primal is extended by at least one assignment



# Primal-dual (c'od)

d) Dual change routine:



$$\delta = \min \{ P_{ij} \mid (i,j) \in -, P_{ij} > 0 \}.$$

The update for dual variables is as on p. 7, except we take the  $\min P_{ij}$  among all labelled rows (no cell on labelled rows becomes infeasible for dual pb.)

- Result
- # of  $(P_{ij} = 0)$  cells increases by 1 at least (primal assignment does not necessarily increase by 1 pair!)
  - $O(n^2)$  iterations