

Computer Science 1820

Solutions for Recommended Exercises

Section 1.3

6. (a) “ \exists ” means “there exists,” so $\exists xN(x)$ means that there exists some member x of the population for which $N(x)$ is true. In English, this quantification is “Some student in my school has visited North Dakota.”
- (b) “ \forall ” means “for all,” so $\forall xN(x)$ means that $N(x)$ is true for every member x of the population. In English, this quantification is “Every student in my school has visited North Dakota.”
- (c) “It is not the case that some student in my school has visited North Dakota.”
- (d) “Some student in my school has not visited North Dakota.”
- (e) “It is not the case that every student in my school has visited North Dakota.”
- (f) “Every student in my school has not visited North Dakota.”
8. (a) “Every rabbit hops.”
- (b) “Every animal is a rabbit and hops,” or “Every animal is a hopping rabbit.”
- (c) “Some rabbit hops.”
- (d) “Some animal is a rabbit and hops,” or “Some animal is a hopping rabbit.”
10. (a) $\exists x(C(x) \wedge D(x) \wedge F(x))$ (b) $\forall x(C(x) \vee D(x) \vee F(x))$
- (c) $\exists x(C(x) \wedge F(x) \wedge (\neg D(x)))$ (d) $\neg \exists x(C(x) \wedge D(x) \wedge F(x))$
- (e) $(\exists xC(x)) \wedge (\exists xD(x)) \wedge (\exists xC(x))$, or $(\exists xC(x)) \wedge (\exists yD(y)) \wedge (\exists zC(z))$

12. (a) $(0) + 1 = 1$, $2(0) = 0$, and $1 > 0$, so $Q(0)$ is true.
- (b) $(-1) + 1 = 0$, $2(-1) = -2$, and $0 > -2$, so $Q(-1)$ is true.
- (c) $(1) + 1 = 2$, $2(1) = 2$, but $2 \not> 2$, so $Q(1)$ is false.
- (d) There *does* exist an integer x for which $Q(x)$ is true: 0! So, $\exists x Q(x)$ is true.
- (e) $Q(x)$ is *not* true for all integers x ; specifically, $Q(1)$ is false! So, $\forall x Q(x)$ is false.
- (f) There does exist an integer x for which $Q(x)$ is false (1 is such an integer), so there exists an integer for which $\neg Q(x)$ is true. Ergo, $\exists x \neg Q(x)$ is true.
- (g) $Q(x)$ is not false for every integer x ($Q(0)$ is not false), so $\neg Q(x)$ is not true for every integer x . Hence, $\forall x \neg Q(x)$ is false.
14. (a) $(-1)^3 = -1$, so there exists a real number x for which $x^3 = -1$ is true; thus, the given statement is true.
- (b) $(0.5)^4 = 0.0625$, $(0.5)^2 = 0.25$, and $0.0625 < 0.25$, so $x = 0.5$ makes $x^4 < x^2$ true; therefore, the given statement is true.
- (c) Let x be any real number. Then, $(-x)^2 = (-1)^2 x^2 = (1)x^2 = x^2$, so $(-x)^2 = x^2$ is true for every real number. Consequently, the given statement is true.
- (d) $2(-1) = -2$, but $-2 \not> -1$, so $2x > x$ is *not* true for every real number x ; ergo, the given statement is false.
18. (a) $P(-2) \vee P(-1) \vee P(0) \vee P(1) \vee P(2)$
- (b) $P(-2) \wedge P(-1) \wedge P(0) \wedge P(1) \wedge P(2)$
- (c) $(\neg P(-2)) \vee (\neg P(-1)) \vee (\neg P(0)) \vee (\neg P(1)) \vee (\neg P(2))$
- (d) $(\neg P(-2)) \wedge (\neg P(-1)) \wedge (\neg P(0)) \wedge (\neg P(1)) \wedge (\neg P(2))$

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(e) $\neg(P(-2) \vee P(-1) \vee P(0) \vee P(1) \vee P(2))$

Using De Morgan's Law, this becomes: $(\neg P(-2)) \wedge (\neg P(-1)) \wedge (\neg P(0)) \wedge (\neg P(1)) \wedge (\neg P(2))$

(f) $\neq (P(-2) \wedge P(-1) \wedge P(0) \wedge P(1) \wedge P(2))$

Using De Morgan's Law, this becomes: $(\neg P(-2)) \vee (\neg P(-1)) \vee (\neg P(0)) \vee (\neg P(1)) \vee (\neg P(2))$

20. (a) $P(-5) \vee P(-3) \vee P(-1) \vee P(1) \vee P(3) \vee P(5)$

(b) $P(-5) \wedge P(-3) \wedge P(-1) \wedge P(1) \wedge P(3) \wedge P(5)$

(c) We can express the given quantified statement in English as "For every x that is not 1, $P(x)$ is true," which in turn allows us to express it as

$$P(-5) \wedge P(-3) \wedge P(-1) \wedge P(3) \wedge P(5).$$

(d) We can express the given quantified statement in English as "There is some x that is nonnegative for which $P(x)$ is true," which is logically equivalent to $P(1) \vee P(3) \vee P(5)$.

(e) We express the given statement in English as "There is some x for which $P(x)$ is not true, and for every negative x , $P(x)$ is true." Note the consequence of this conjunction: whatever value of x makes $P(x)$ false must be nonnegative! Ergo, the statement becomes

$$P(-5) \wedge P(-3) \wedge P(-1) \wedge [(\neg P(1)) \vee (\neg P(3)) \vee (\neg P(5))].$$

30. (a) $P(1,3) \vee P(2,3) \vee P(3,3)$

(b) $P(1,1) \wedge P(1,2) \wedge P(1,3)$

(c) $(\neg P(2,1)) \vee (\neg P(2,2)) \vee (\neg P(2,3))$

(d) $(\neg P(1,2)) \wedge (\neg P(2,2)) \wedge (\neg P(3,2))$

32. To begin, let the domain in each part of this exercise be the set of all (living) animals. Also, we will use these four logical equivalences:

$$\neg(\forall xP(x)) \equiv \exists x\neg P(x), \quad \neg(\exists xP(x)) \equiv \forall x\neg P(x), \quad \neg(A \rightarrow B) \equiv (A \wedge \neg B), \quad \text{and} \quad \neg(A \wedge B) \equiv (A \rightarrow \neg B).$$

(a) We let $D(x)$ represent “ x is a dog” and let $F(x)$ denote “ x has fleas.” Then, the original statement is $\forall x(D(x) \rightarrow F(x))$. Its negation is therefore

$$\neg[\forall x(D(x) \rightarrow F(x))] \equiv \exists x\neg(D(x) \rightarrow F(x)) \equiv \exists x(D(x) \wedge \neg F(x)).$$

In English, this negation is “Some dog does not have fleas.”

(b) Letting $H(x)$ mean “ x is a horse” and $A(x)$ stand for “ x can add,” the original statement becomes $\exists x(H(x) \wedge A(x))$. Its negation is

$$\neg[\exists x(H(x) \wedge A(x))] \equiv \forall x\neg(H(x) \wedge A(x)) \equiv \forall x(H(x) \rightarrow \neg A(x)),$$

whose English form is “Every horse cannot add.”

(c) Letting $K(x)$ be “ x is a koala” and $C(x)$ be “ x can climb,” the original statement is $\forall x(K(x) \rightarrow C(x))$. Its negation is

$$\neg[\forall x(K(x) \rightarrow C(x))] \equiv \exists x\neg(K(x) \rightarrow C(x)) \equiv \exists x(K(x) \wedge \neg C(x)),$$

whose English form is “Some koala cannot climb.”

(d) Letting $M(x)$ represent “ x is a monkey” and $F(x)$ denote “ x can speak French,” the original statement is $\neg\exists x(M(x) \wedge F(x))$. Its negation is

$$\neg[\neg\exists x(M(x) \wedge F(x))] \equiv \exists x(M(x) \wedge F(x)),$$

whose English form is “Some monkey can speak French.”

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- (e) Letting $P(x)$ be “ x is a pig,” $S(x)$ be “ x can swim,” and $C(x)$ be “ x can catch fish,” the original statement is $\exists x(P(x) \wedge (S(x) \wedge C(x)))$. Its negation is

$$\begin{aligned}\neg[\exists x(P(x) \wedge (S(x) \wedge C(x)))] &\equiv \forall x\neg(P(x) \wedge (S(x) \wedge C(x))) \\ &\equiv \forall x(P(x) \rightarrow \neg(S(x) \wedge C(x))) \equiv \forall x(P(x) \rightarrow (\neg S(x) \vee \neg C(x))),\end{aligned}$$

whose English form is “Every pig cannot swim or cannot catch fish.”

34. (a) Let the domain be the set of all drivers, and let $O(x)$ represent “ x obeys the speed limit.” Then, the original statement is $\exists x\neg O(x)$, and its negation is

$$\neg[\exists x\neg O(x)] \equiv \forall x\neg(\neg O(x)) \equiv \forall xO(x).$$

In English, this negation is “Every driver obeys the speed limit.”

- (b) Let the domain be the set of all movies, $W(x)$ denote “ x is Swedish,” and $S(x)$ mean “ x is serious.” Then, the original statement is $\forall x(W(x) \rightarrow S(x))$, so its negation is

$$\neg[\forall x(W(x) \rightarrow S(x))] \equiv \exists x\neg(W(x) \rightarrow S(x)) \equiv \exists x(W(x) \wedge \neg S(x)).$$

In English, this negation is “Some Swedish movie is not serious.”

- (c) Letting the domain be the set of all (living) persons and $K(x)$ represent “ x can keep a secret,” the original statement becomes $\neg\exists xK(x)$. Then, its negation is

$$\neg[\neg\exists xK(x)] \equiv \exists xK(x),$$

whose English form is “Someone can keep a secret.”

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- (d) Let the domain be the set of people in this class, and let $G(x)$ be “ x has a good attitude.” Then, the original statement is $\exists x\neg G(x)$, so its negation is

$$\neg[\exists x\neg G(x)] \equiv \forall x\neg(\neg G(x)) \equiv \forall xG(x).$$

In English, this negation is “Every student in this class has a good attitude.”

36. (a) 0 or 1

(b) $-\sqrt{2}$ or $\sqrt{2}$

(c) 0